

Büchi Objectives in Countable MDPs

Stefan Kiefer, Richard Mayr, Patrick Totzke

Mahsa Shirmohammadi

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Following his departure from Circe's island home of Aeaëa, **Odysseus** braces for the many challenges he will encounter on his journey home to his beloved Ithaca



Dilemma: Between a rock and a hard place

1

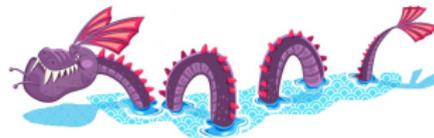


Scylla

strait of Messina



too close
inescapable!



Charybdis

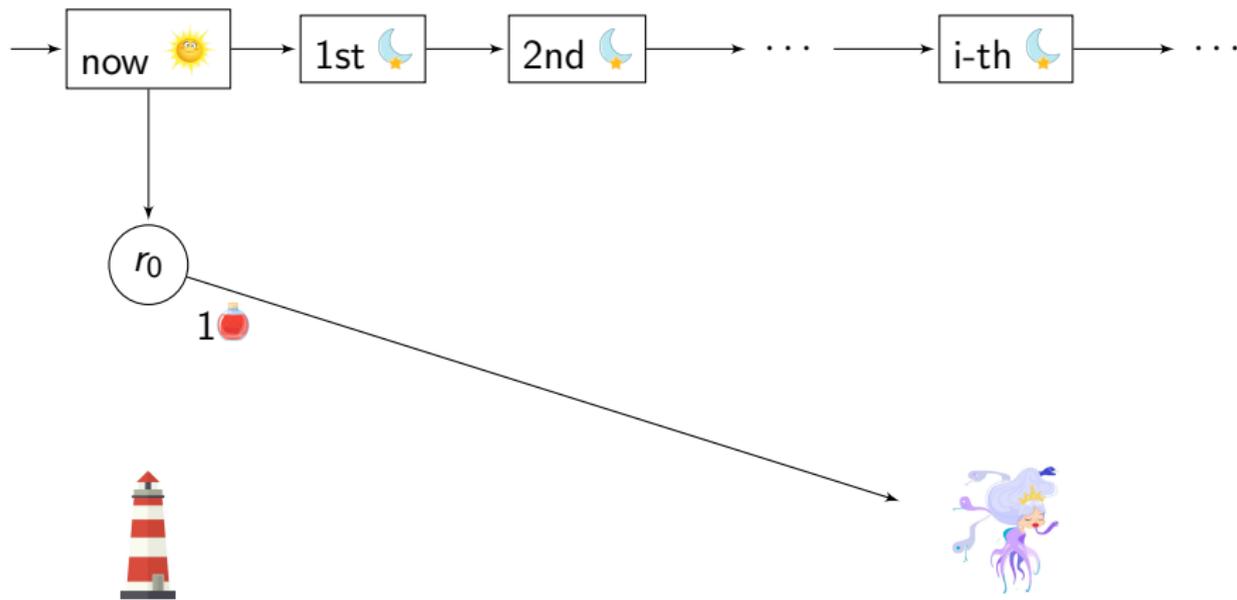
MDP of Odysseus's dilemma

2



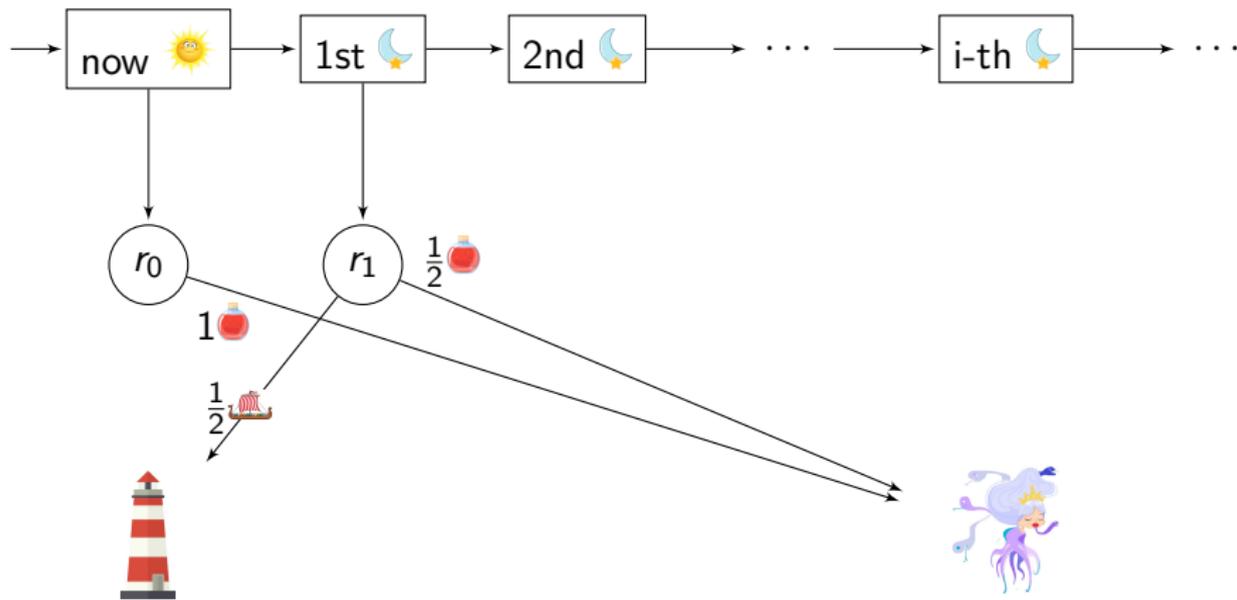
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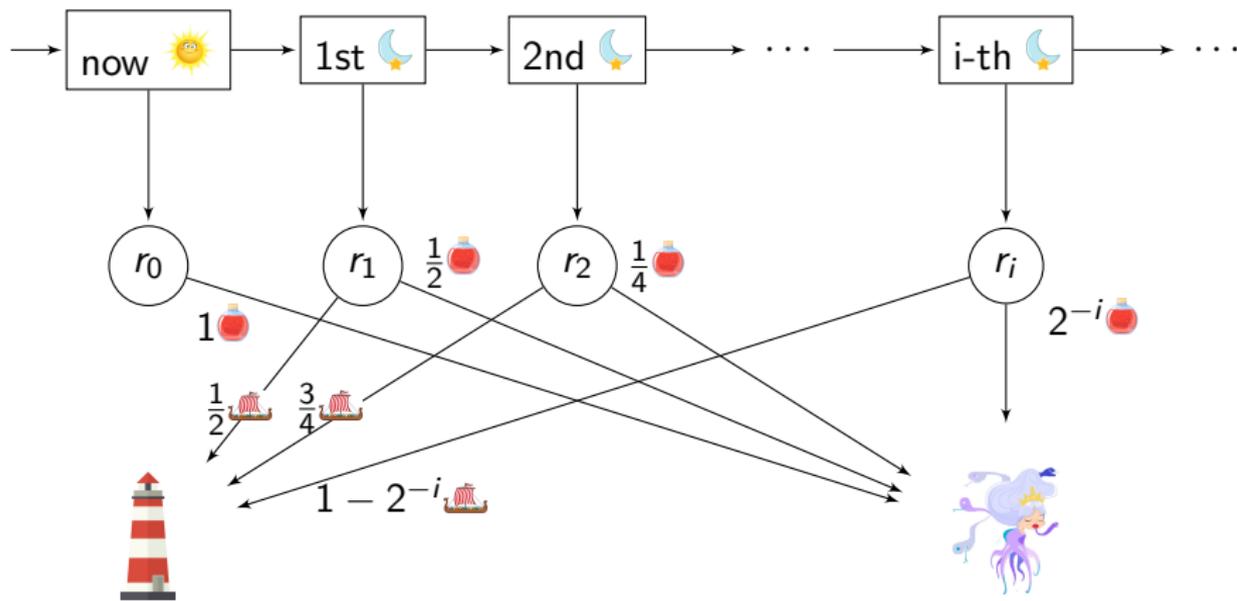
MDP of Odysseus's dilemma

2



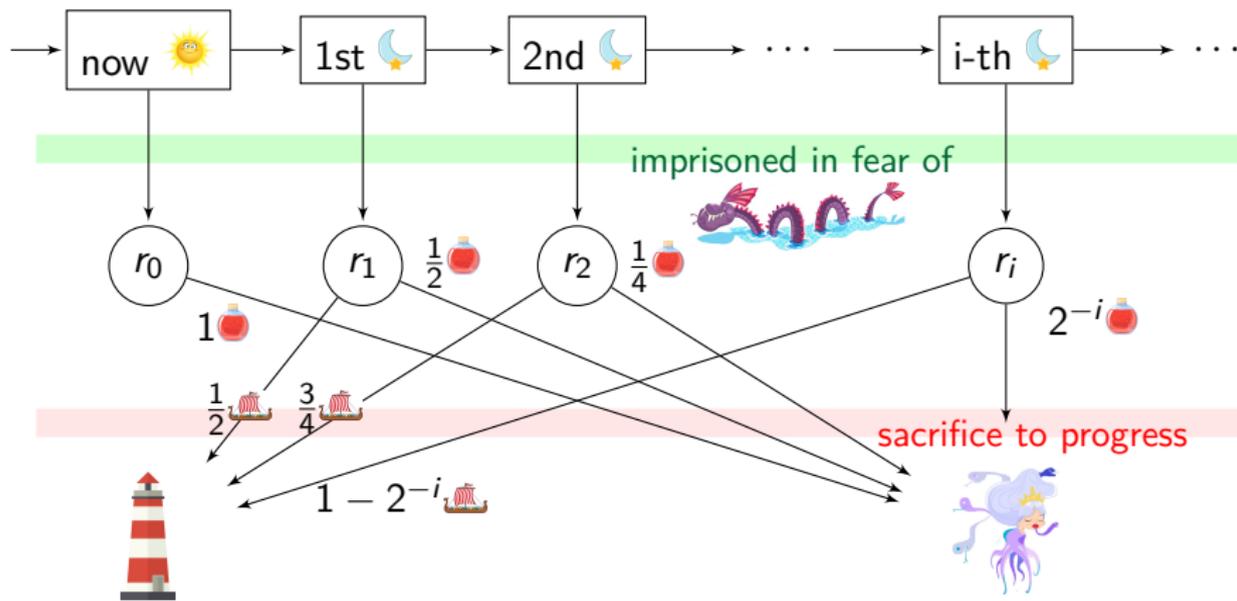
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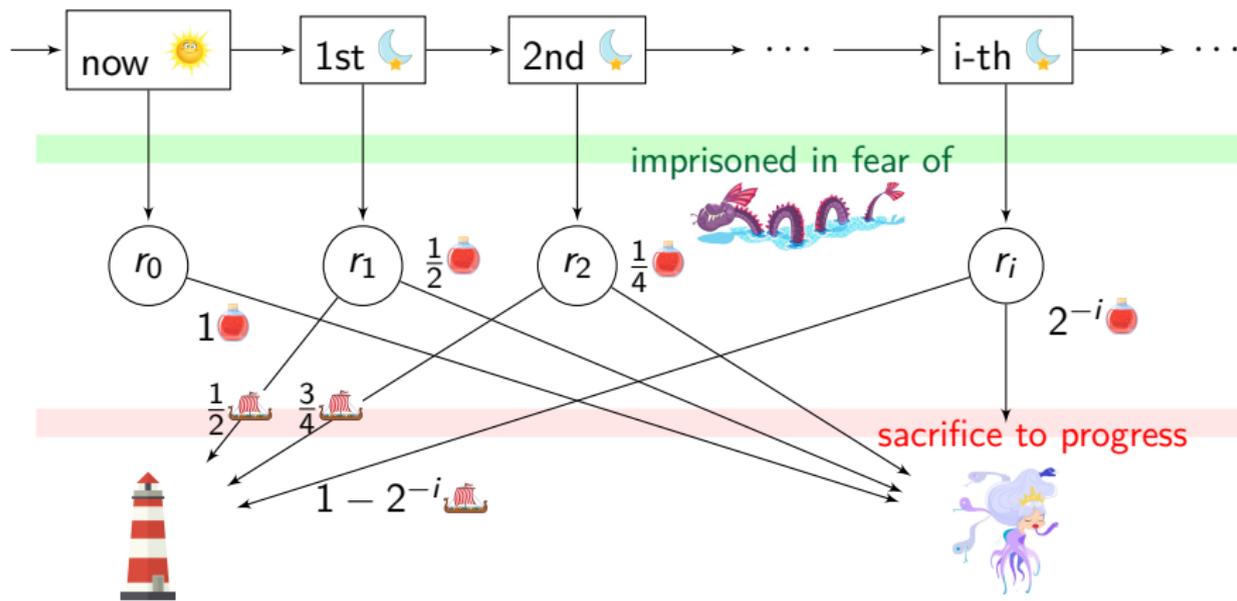
2



MDP of Odysseus's dilemma

2





The value of $\text{Reach}(\text{Lighthouse})$ is 1:

for all $\epsilon > 0$,  manages to get his crew back home with $1 - \epsilon$ probability!

In a recent version:



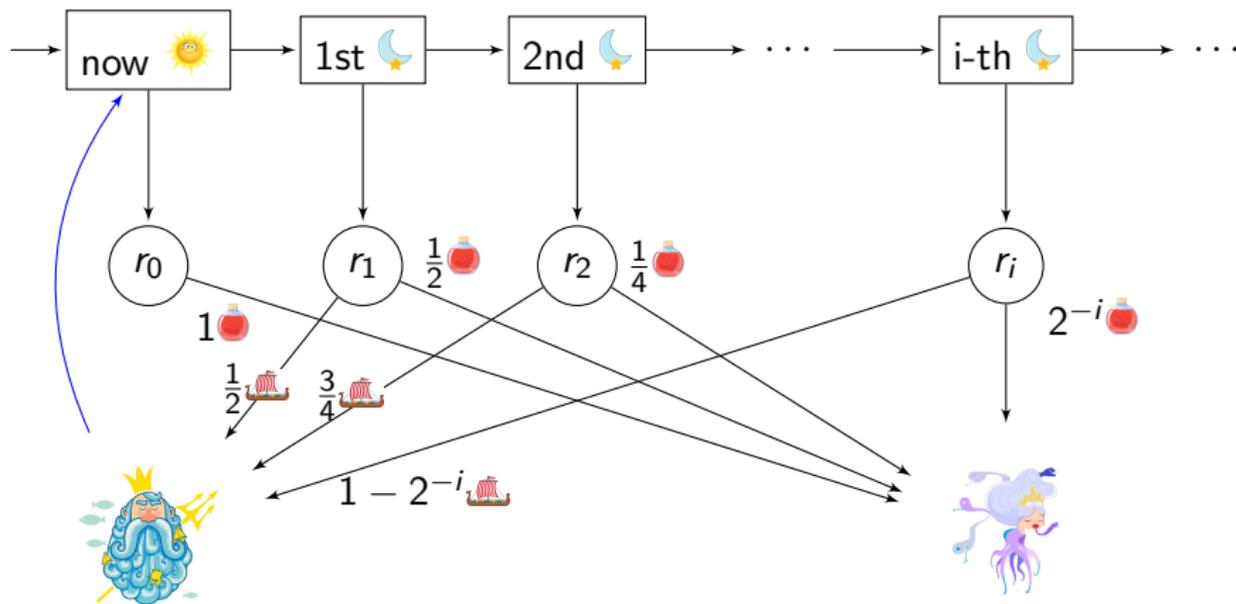
after passing through Scylla meets



Poseidon

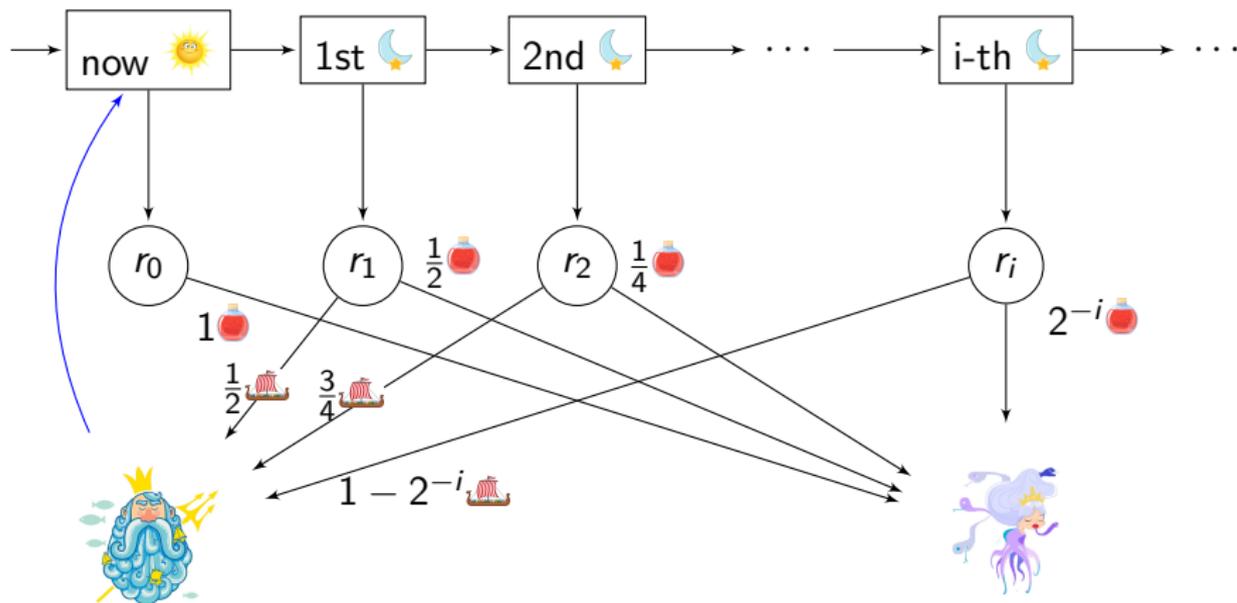
!!! Oops!

... still furious at Odysseus making his son blind.



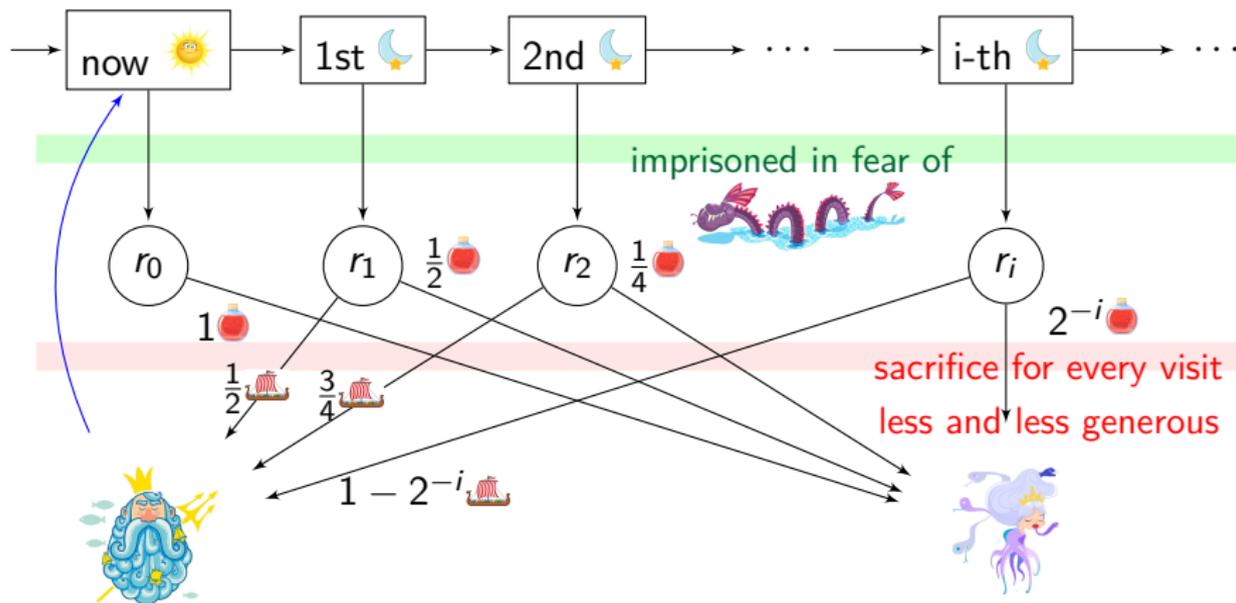
: I make you suffer!

Visit me (i.e., sacrifice to Scylla) **over and over!**



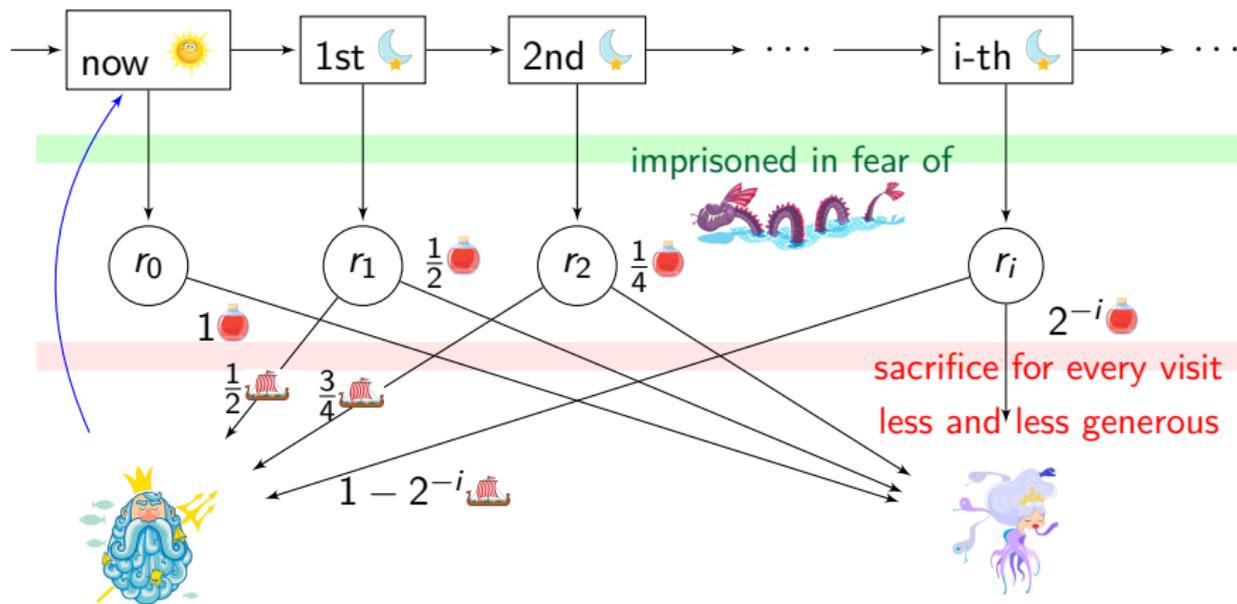
⚡ Is the value of Büchi() one?

for all $\epsilon > 0$, can  visit  ∞ -times with probability at least $1 - \epsilon$?



⚡ Is the value of Büchi() one? **YES!**

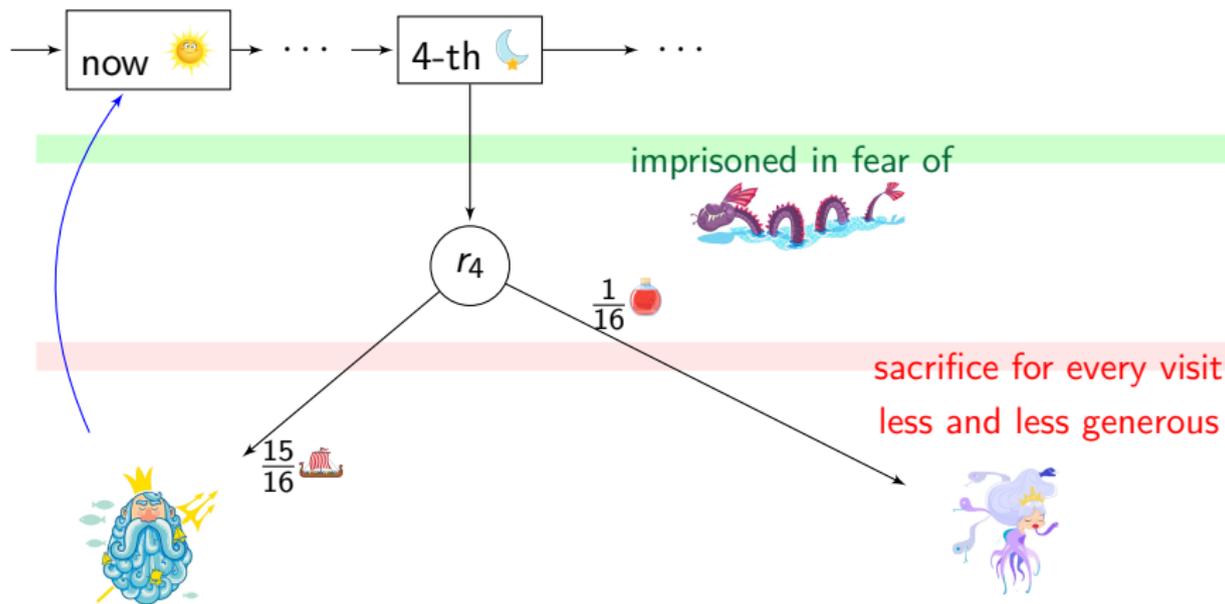
for all $\epsilon > 0$, can  visit  ∞ -times with probability at least $1 - \epsilon$?



⚡ Let $\epsilon = \frac{1}{8}$, let's see how $\frac{7}{8}$ of crew visit  ∞ -times.

How? 1st visit

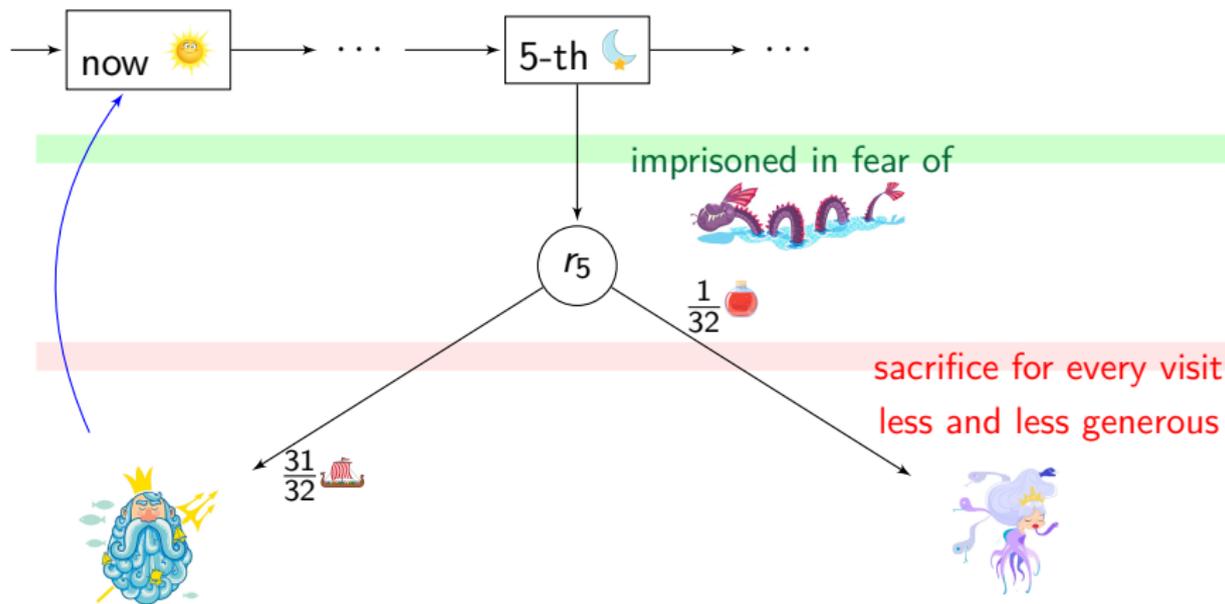
5



⚡ Let $\epsilon = \frac{1}{8}$, let's see how $\frac{7}{8}$ of crew visit  ∞ -times.
the total sacrifice so far: $\frac{1}{16}$

How? 2nd visit

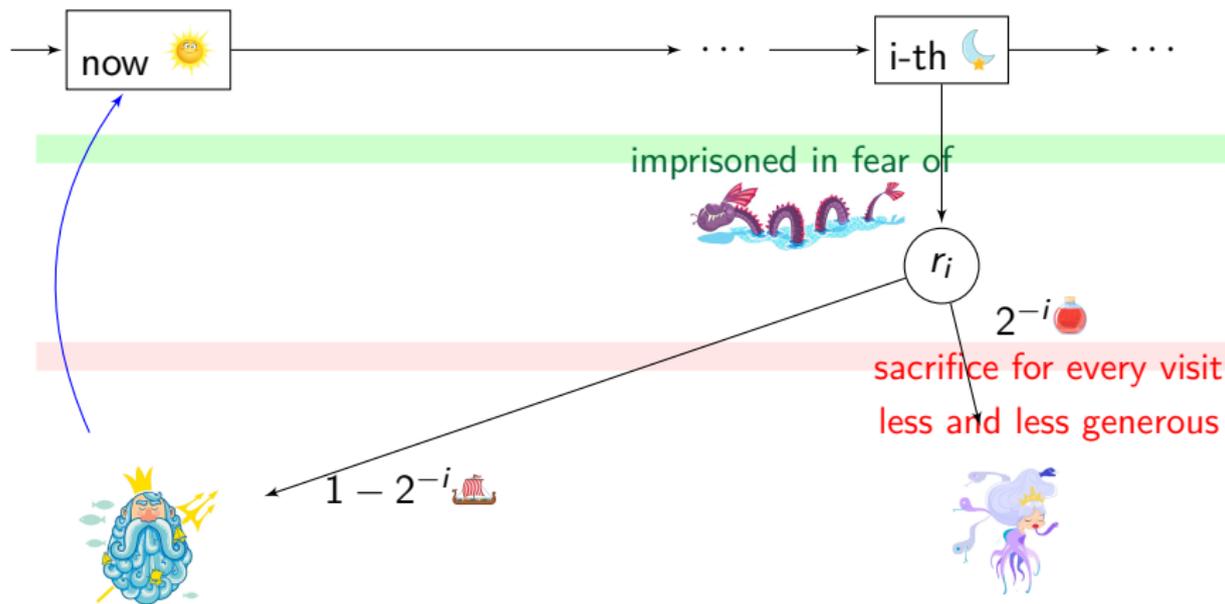
5



⚡ Let $\epsilon = \frac{1}{8}$, let's see how $\frac{7}{8}$ of crew visit  ∞ -times.
the total sacrifice so far: $\frac{1}{16} + \frac{1}{32}$

How? being strategic!

5

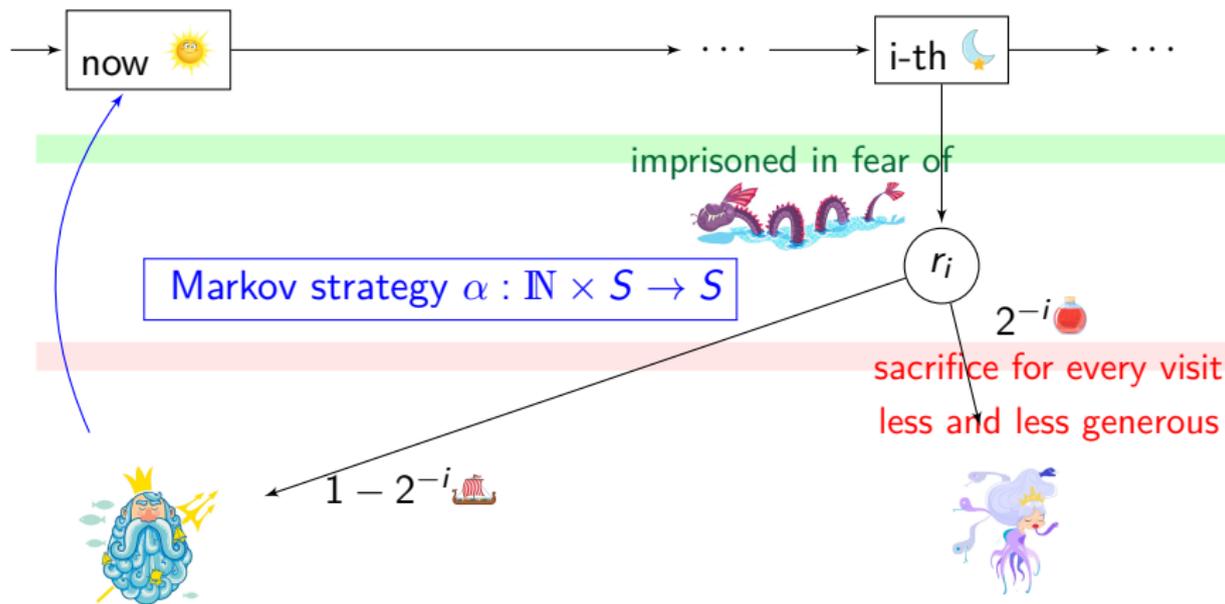


⚡ Let $\epsilon = \frac{1}{8}$, let's see how $\frac{7}{8}$ of crew visit  ∞ -times.

the total sacrifice so far: $\frac{1}{16} + \frac{1}{32} + \dots + \frac{1}{2^i} + \dots = \sum_{i=4}^{\infty} \frac{1}{2^i} = \frac{1}{8}$

How? Markov Strategy

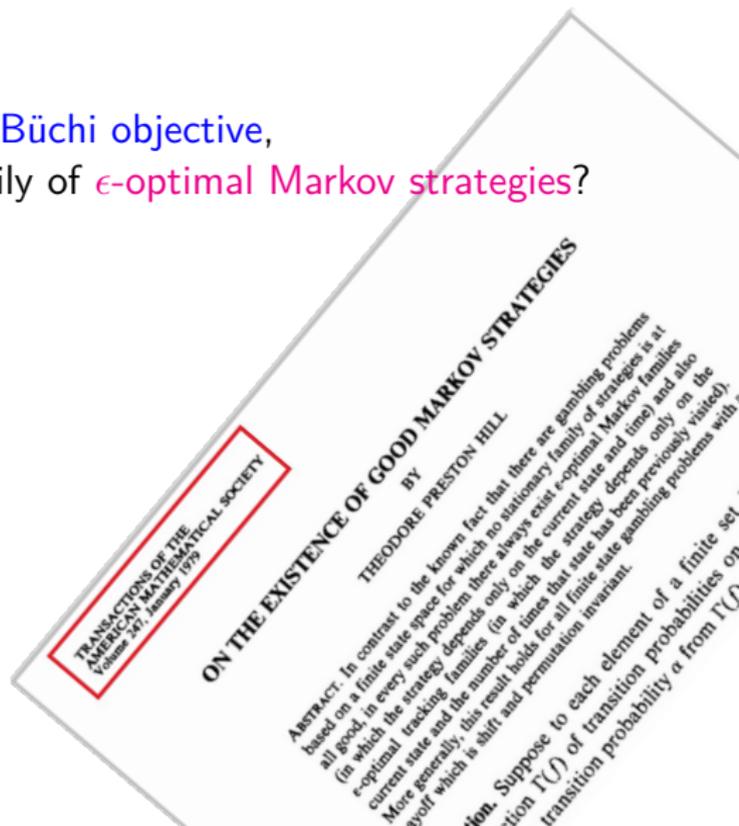
6



⚡ Let $\epsilon = \frac{1}{8}$, let's see how $\frac{7}{8}$ of crew visit  ∞ -times.
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does there always exist a family of **ϵ -optimal Markov strategies**?

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Q1. (cf. [Hi]) Do good Markov strategies exist in all countable-state goal problems with objective of hitting the goal infinitely often?

IEEE Technol
4-132, 1998

TRANSACTIONS OF THE
AMERICAN MATHEMATICAL SOCIETY
347, January 1999

THE EXISTENCE OF GOOD MARKOV STRATEGIES

BY
THEODORE PRESTON HILL

As contrast to the known fact that there are gambling problems in every such space for which no stationary family of strategies is at all good, every such problem there always exist ϵ -optimal Markov families of current state and the number of times that state has been previously visited. More generally, this result holds for all finite state gambling problems with payoff which is shift and permutation invariant.

Suppose to each element of a finite set $\Gamma(\mathcal{U})$ of transition probabilities on transition probability α from $\Gamma(\mathcal{U})$

For countably infinite MDPs and Büchi objective, does there always exist a family of ϵ -optimal Markov strategies?

Q1. (cf. [Hi]) Do good Markov strategies exist in all countable-state goal problems with objective of hitting the goal infinitely often?

⚡ we answer this question :)

Goal Problems in Gambling Theory*

Theodore P. Hill
Center for Applied Probability and
School of Mathematics
Georgia Institute of Technology
Atlanta, GA 30332-0160

Wisc. Technol.
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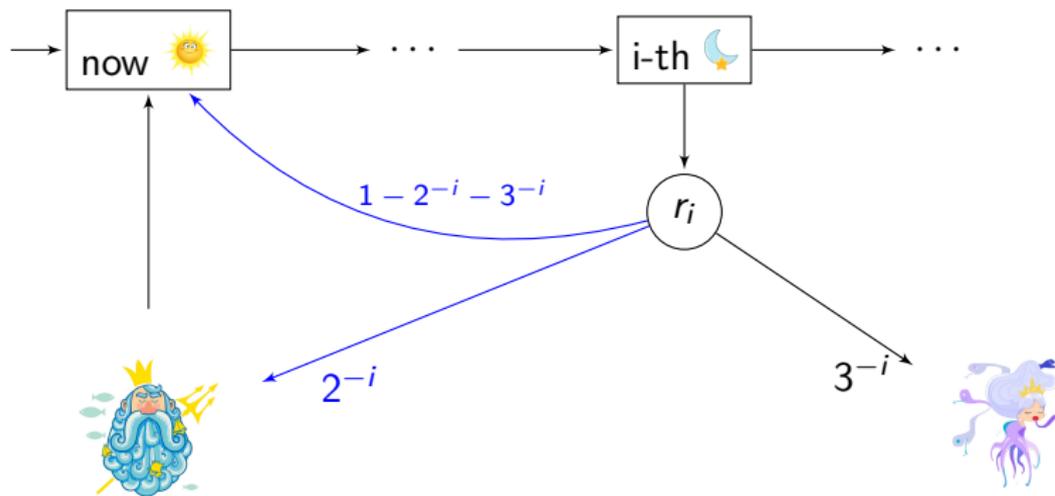
For **countably infinite MDPs** and **Büchi objective**,
does there always exist a family of **ϵ -optimal Markov strategies**?

▷ Is it all about **reducing the risk** of facing dangerous monsters?



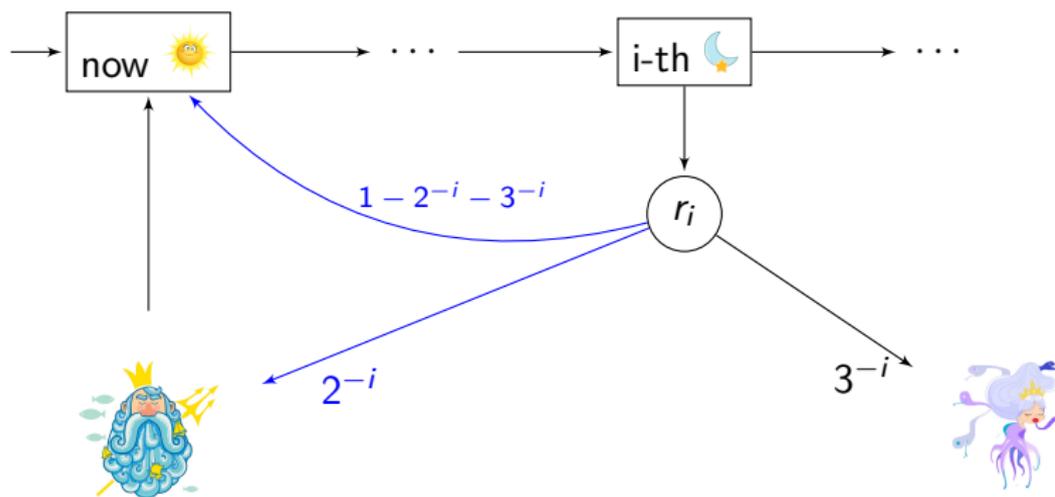
2-nd challenge

8



2-nd challenge

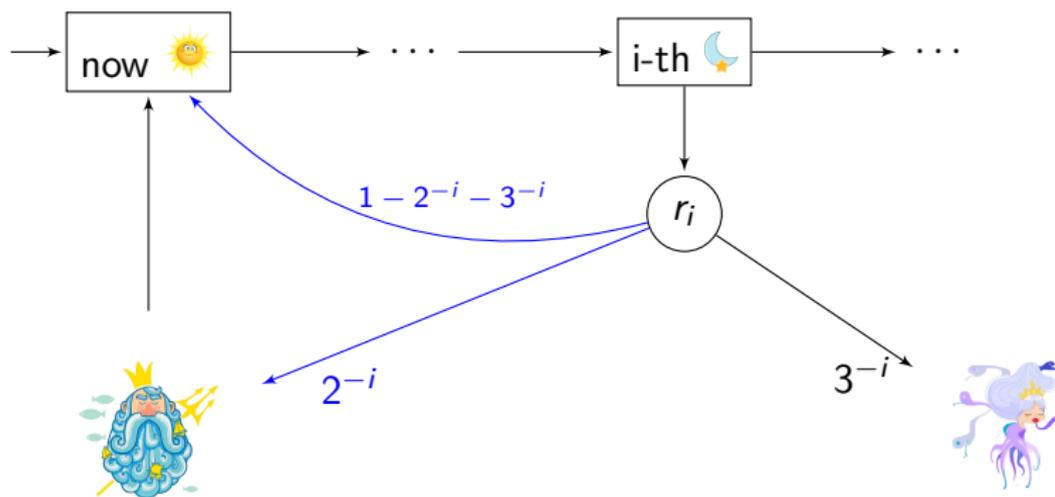
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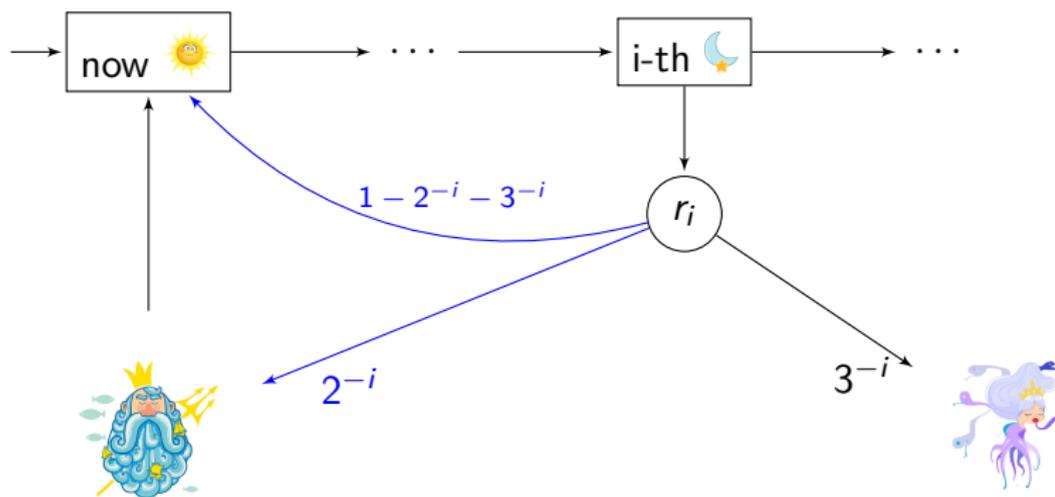
▷ The Markov strategy that, after i -th visit to ☀️, picks r_{i+1} attains 0!

2-nd challenge

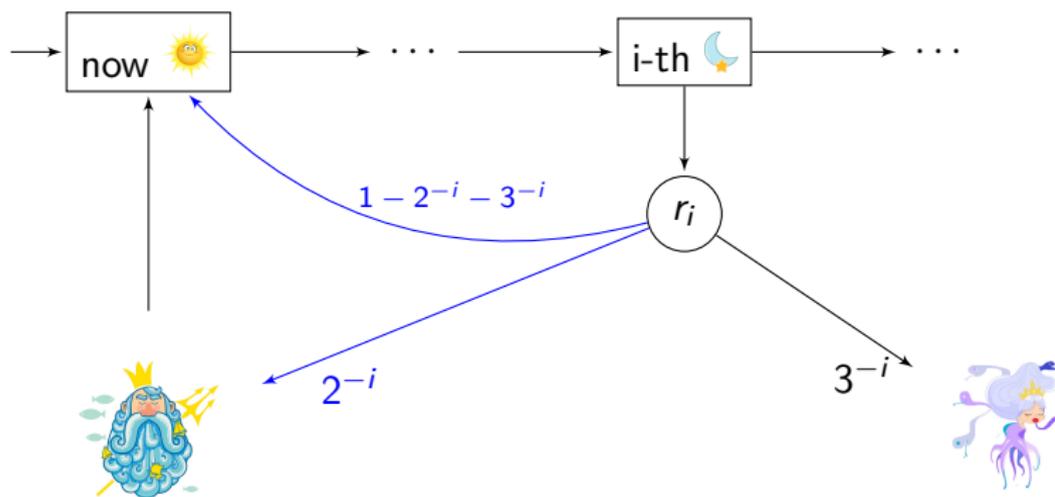
8



- ▷ The Markov strategy that, after i -th visit to ☀, picks r_{i+1} attains 0!
the expected number of visits to Poseidon is at most 1
 $< \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots = 1$



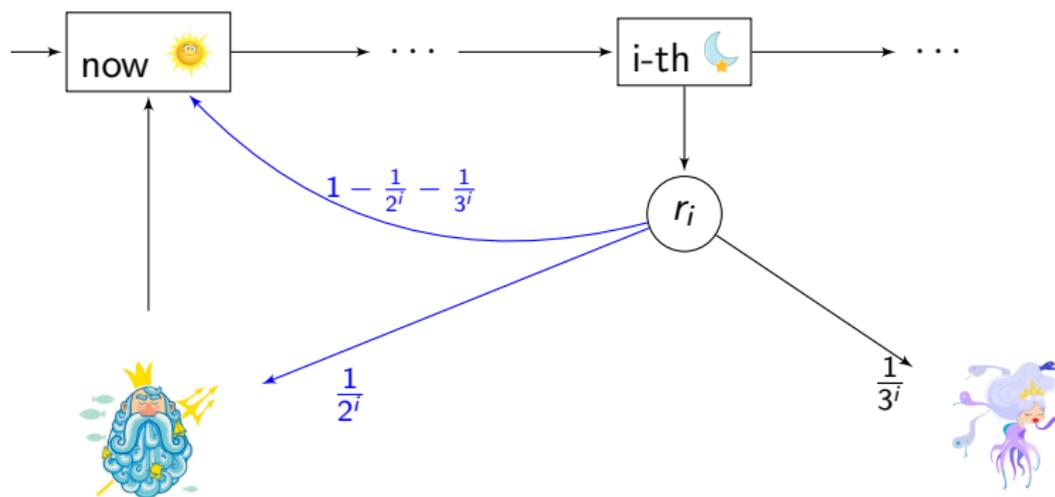
▷ A strategy that picks each r_i for 2^i times achieves Büchi **positively!**



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Bound the total sacrifice by $1 - c$ (technical).

The probability of revisit  after each visit $\geq c$



▷ A strategy that picks each r_i for 2^i times achieves Büchi **positively!**

What is the probability to not visiting Poseidon after i -th phase (for large i)

$$\approx \prod_{k=i}^{\infty} c(1 - \frac{1}{2^k})^{2^k} = 0$$

(since $\sum_{k=i}^{\infty} 2^k \log(c(1 - \frac{1}{2^k}))$ is non-convergent)

For **countably infinite MDPs** and **Büchi objective**,
does there always exist a family of **ϵ -optimal Markov strategies**?

- ▷ it is not all about **reducing the risk** of facing dangerous monsters
- ▷ but rather about a good **compromise** between progress and loss



For **countably infinite MDPs** and **Büchi objective**,
does there always exist a family of ϵ -optimal Markov strategies?

NOOOOoOO!

- ▷ it is not all about **reducing the risk** of facing dangerous monsters
- ▷ but rather about a good **compromise** between progress and loss



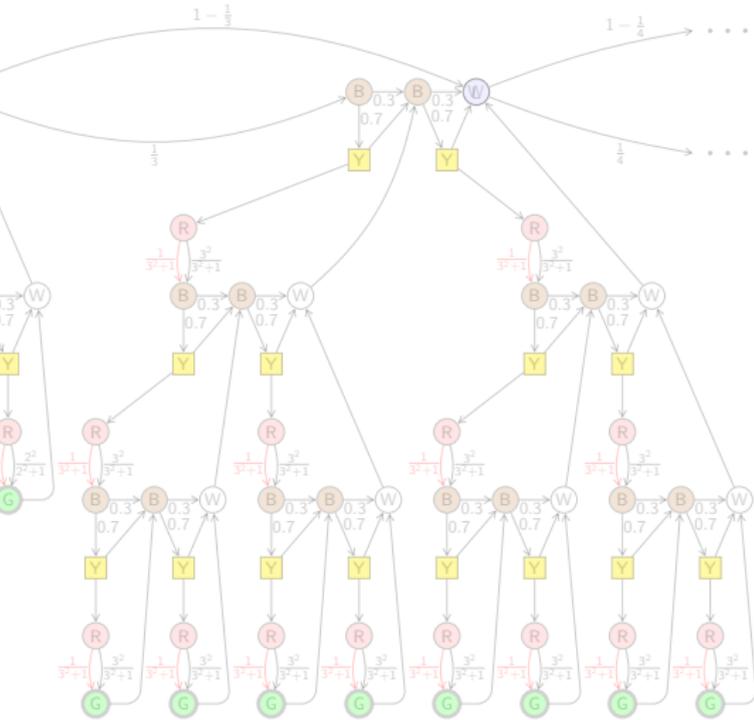
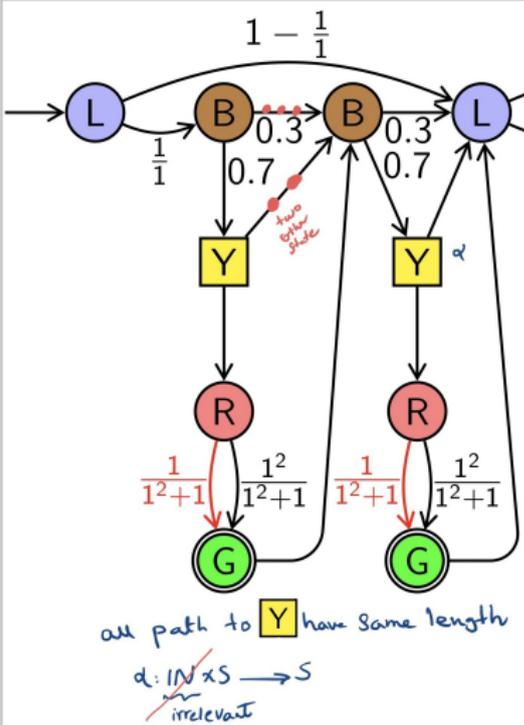
For **countably infinite MDPs** and **Büchi objective**,
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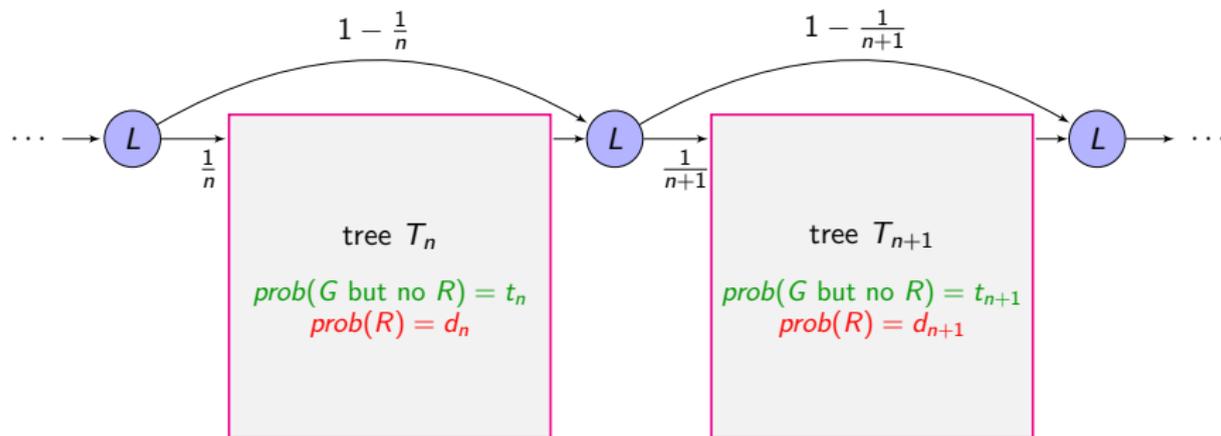
▷ We build an acyclic MDP where **ϵ -optimal strategies** cannot be Markov.

Markov strategy $\alpha : \mathbb{N} \times S \rightarrow S$

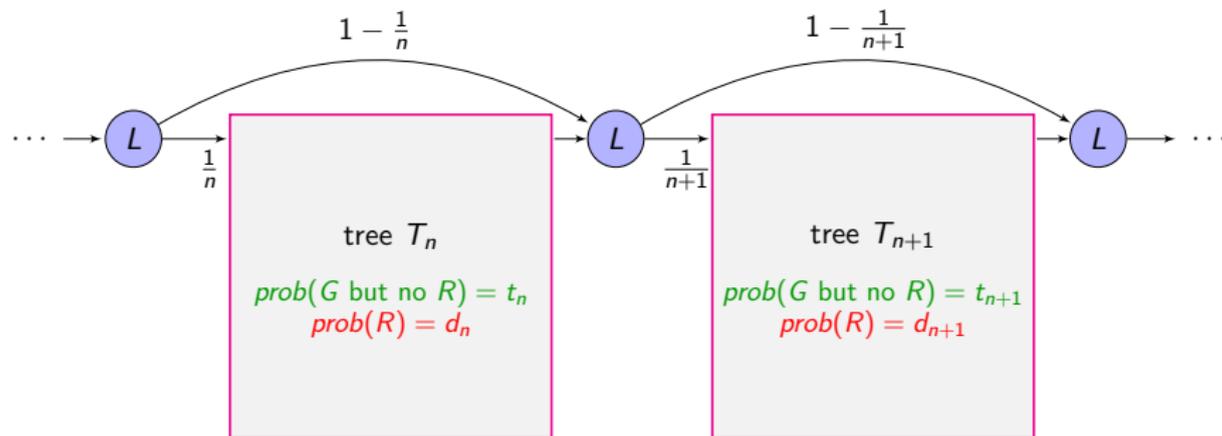
Counter-example



Claim. For Büchi(G) and no R-edge, all Markov strategies attain only 0!



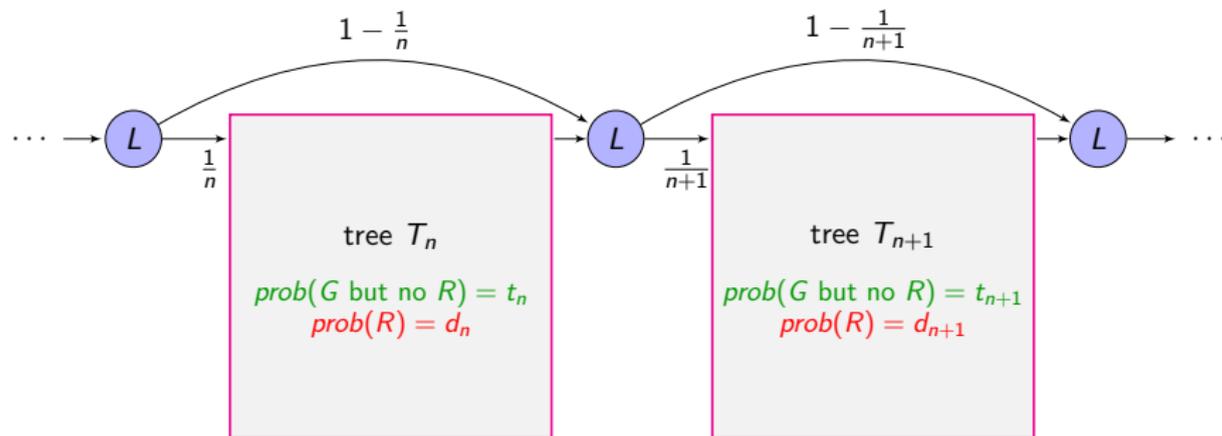
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Expected number of visits to G is $\sum \frac{1}{n} t_n$

$\sum \frac{1}{n} t_n$ must be divergent!

Claim. For Büchi(G) and no R-edge, all Markov strategies attain only 0!



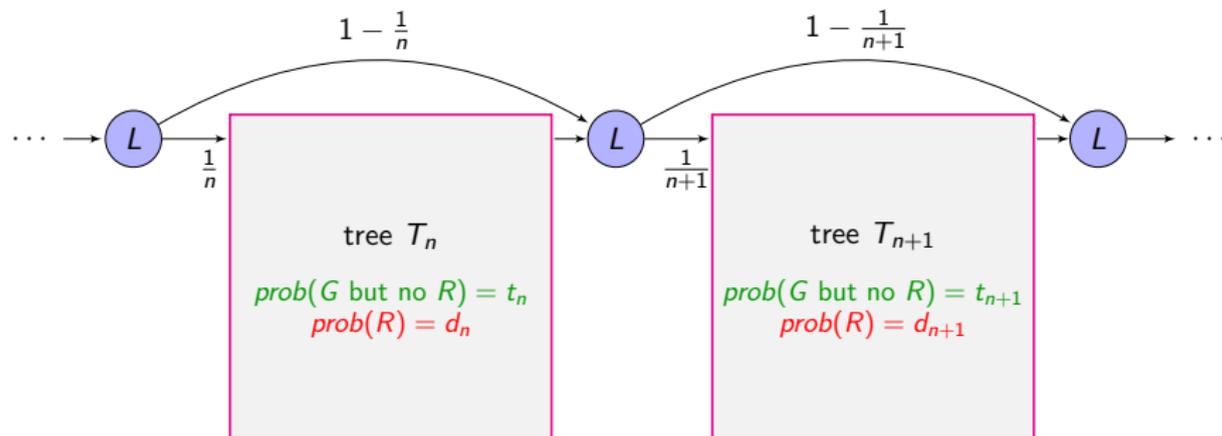
Expected number of visits to G is $\sum \frac{1}{n} t_n$

The probability of R is $\leq \sum \frac{1}{n} d_n$

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The probability of R is $\leq \sum \frac{1}{n} d_n$

$\sum \frac{1}{n} d_n$ must be convergent

By a careful analysis we shows that $d_n \geq 0.008 t_n$ (difficult).

For **countably infinite MDPs** and **Büchi objective**,
does there always exist a family of **ϵ -optimal Markov strategies**?

NOOOOoOO!

- ▷ We showed an acyclic MDPs that **ϵ -optimal strategies** cannot be Markov; however, the value of **Büchi(G)** is 1 (technical).

For countably infinite MDPs and Büchi objective,
does there always exist a family of ϵ -optimal Markov strategies?

NOOOOOoOO!

Theorem. For Büchi, there are always ϵ -optimal 1-bit Markov strategies.

$\alpha : \mathbb{N} \times S \times \{0, 1\} \rightarrow S$ (necessary and sufficient)

- ▷ For Büchi and acyclic MDPs, there are always ϵ -optimal 1-bit strategies

$$\text{1-bit strategy } \alpha : S \times \{ \text{fox}, \text{rabbit} \} \rightarrow S$$

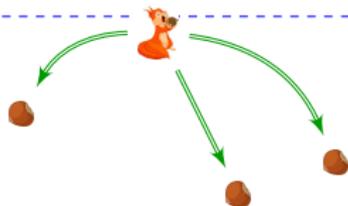
Use 1-bit of memory?

- ▷ For Büchi and acyclic MDPs, there are always ϵ -optimal 1-bit strategies

1-bit strategy $\alpha : S \times \{ \text{fox}, \text{rabbit} \} \rightarrow S$

Fix $\epsilon > 0$.

phase 1



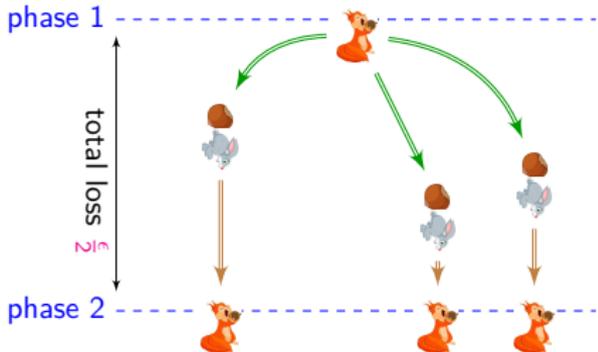
follows $\frac{\epsilon}{4}$ -optimal $\text{Reach}(\bullet)$

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Fix $\epsilon > 0$.



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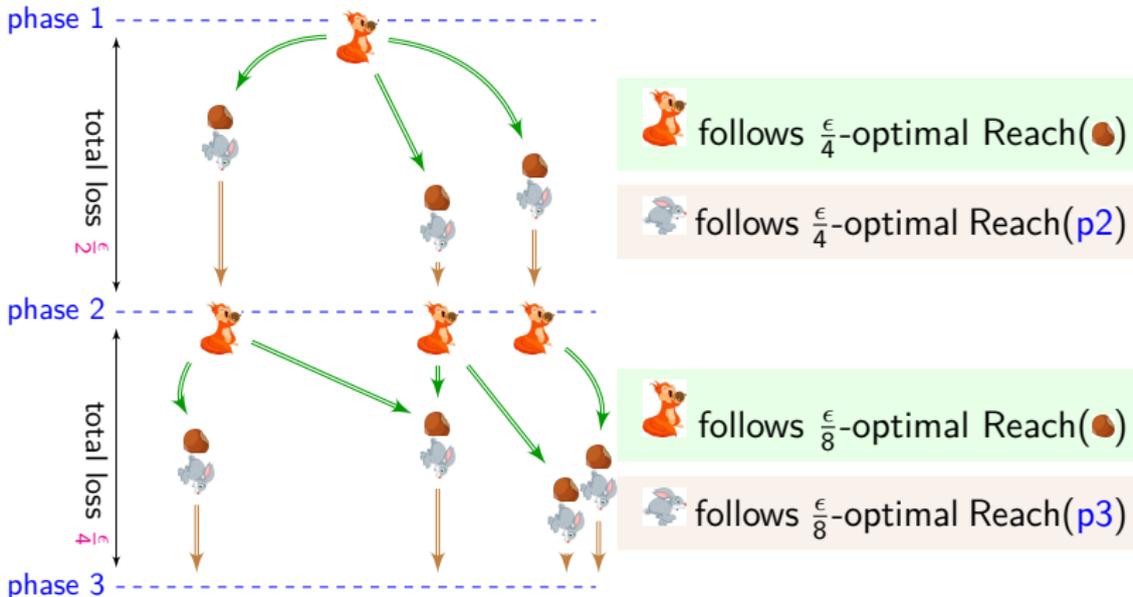
follows $\frac{\epsilon}{4}$ -optimal Reach(p2)

Use 1-bit of memory?

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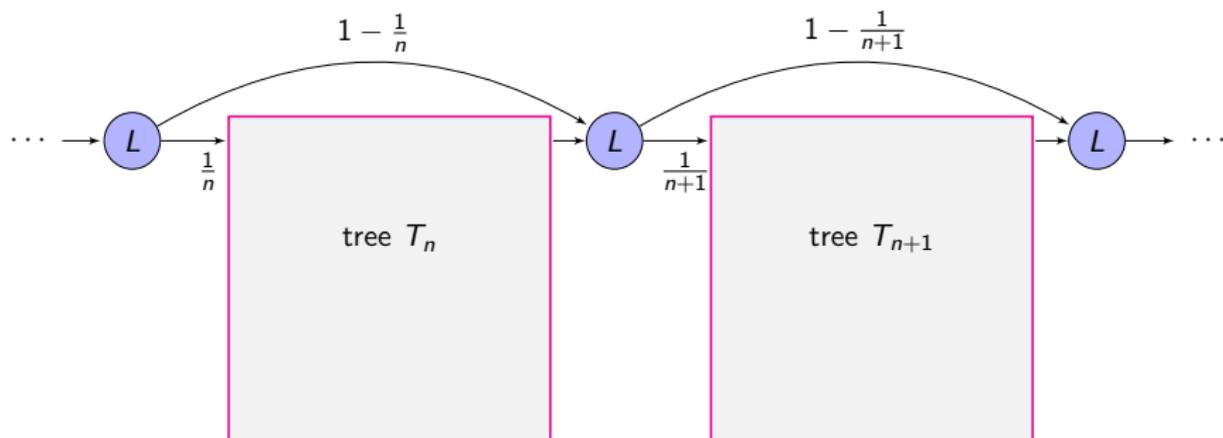
1-bit strategy $\alpha : S \times \{\text{fox}, \text{rabbit}\} \rightarrow S$

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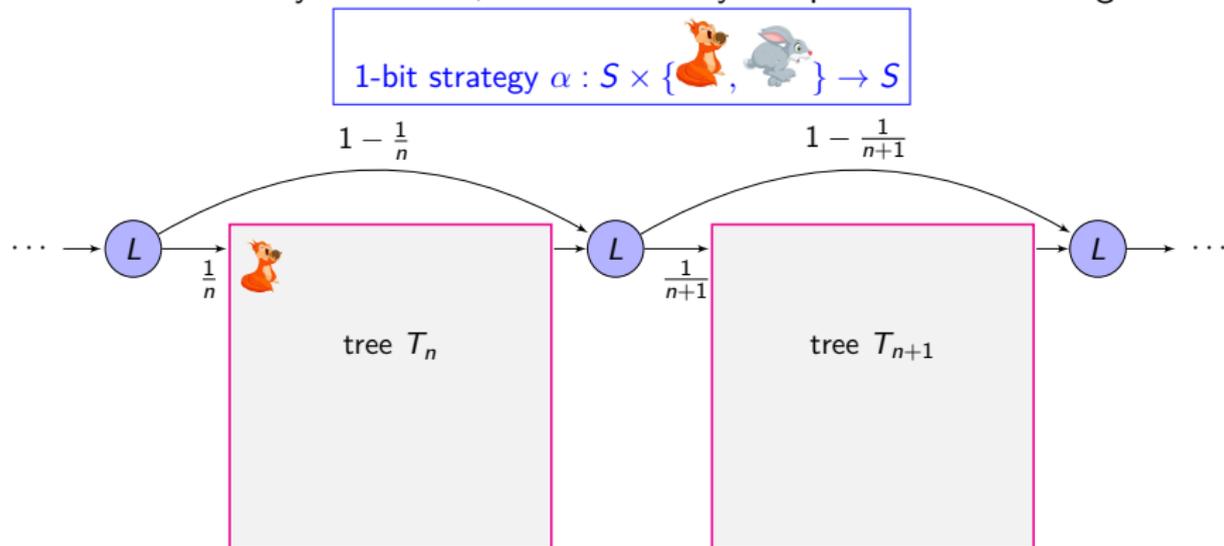
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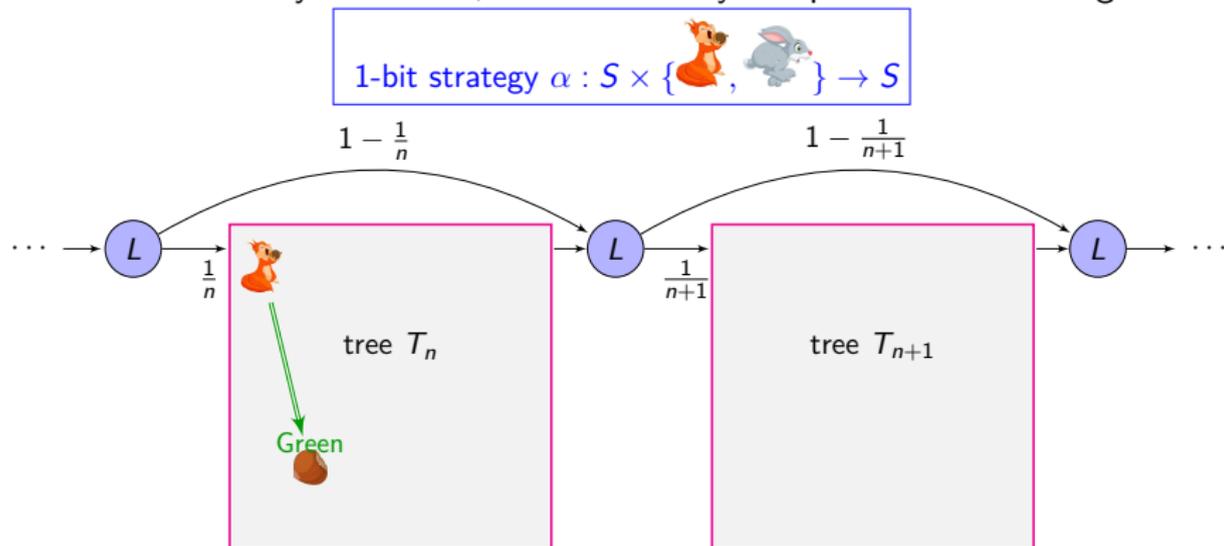
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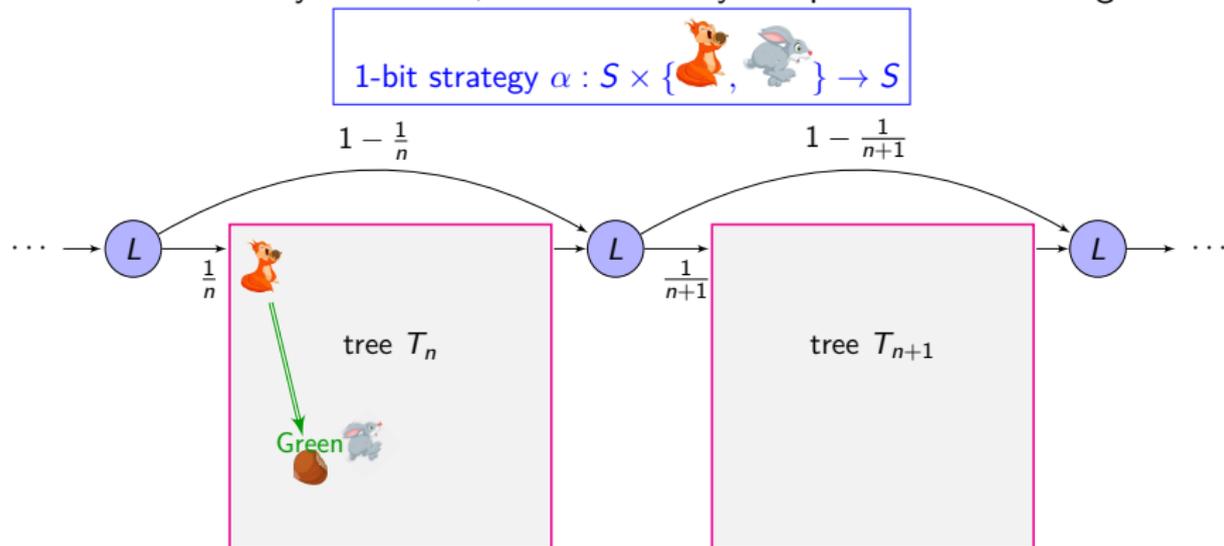
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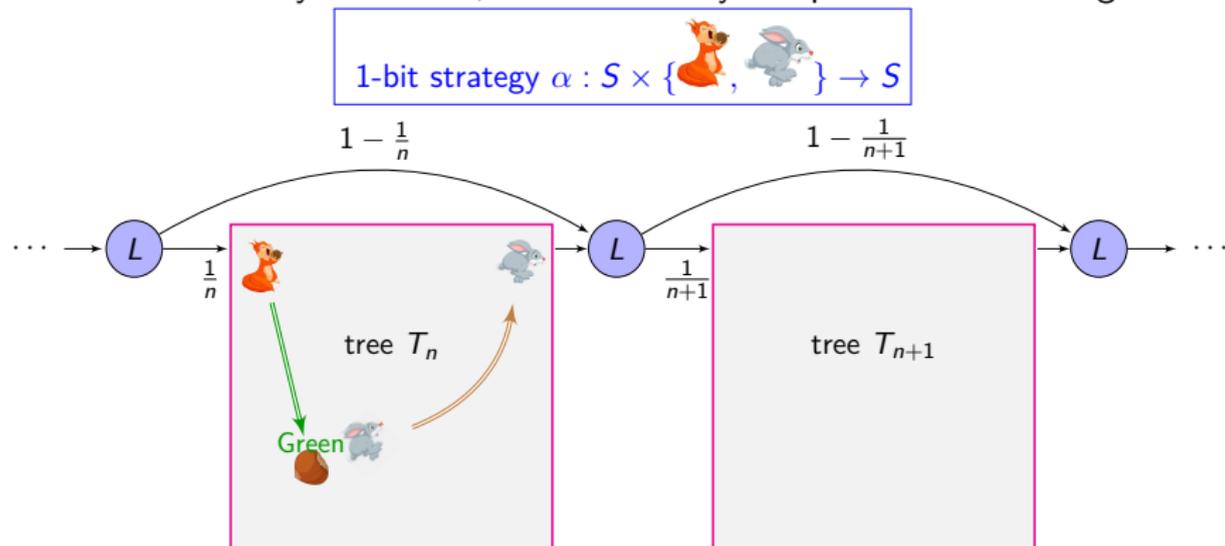
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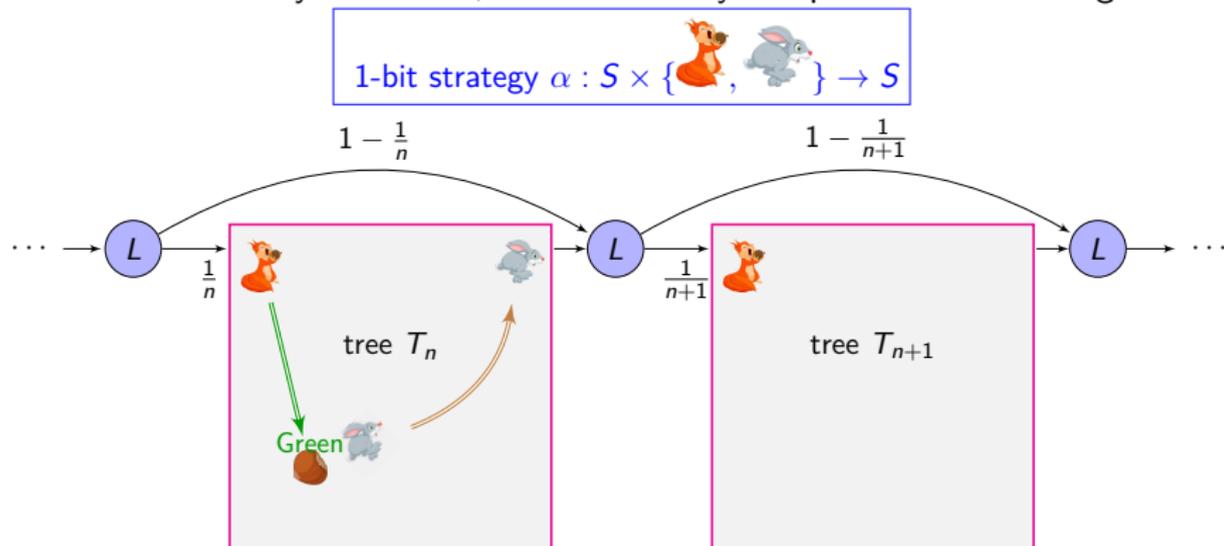
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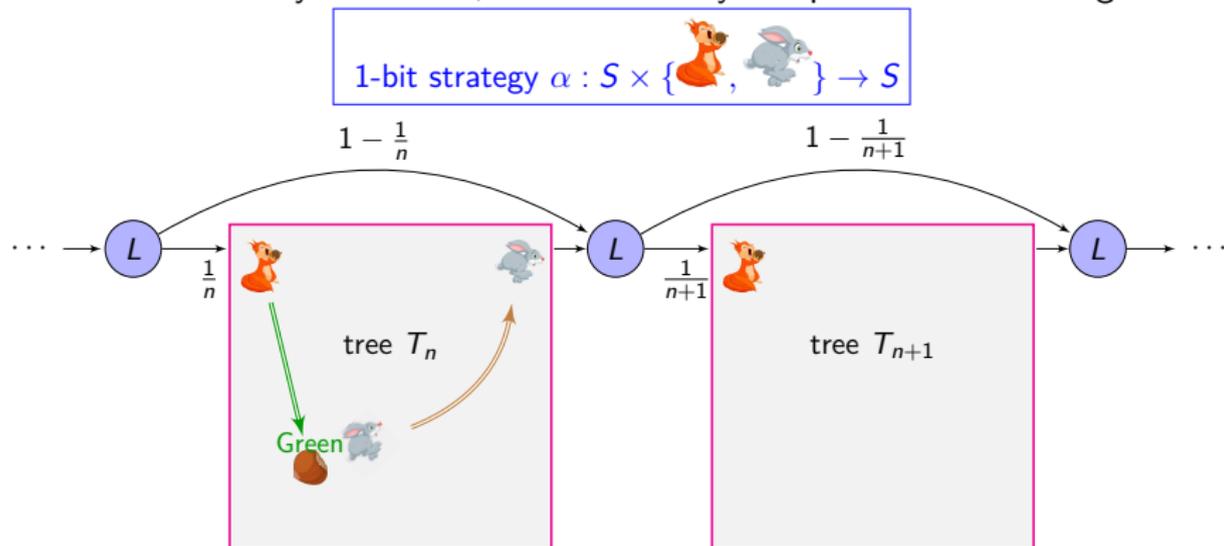
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- ▷ For Büchi and acyclic MDPs, there are always ϵ -optimal 1-bit strategies



For all $\epsilon > 0$, a **starving-squirrel-and-panic-rabbit** strategy achieves $1 - \epsilon$.

Summary: Strategy complexity

*

For countably infinite MDPs and Büchi objective, does there always exist a family of ϵ -optimal Markov strategies?

counter-example.

Theorem. 1-bit Markov strategies are necessary and sufficient.

Goal Problems in Gambling Theory*

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Center for Applied Probability and
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6(2):125-132, 1989

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ON THE EXISTENCE OF GOOD MARKOV STRATEGIES
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ABSTRACT. In contrast to the known fact that there are gambling problems on a finite state space for which no stationary family of strategies is at all ϵ -optimal, every such problem there always exist ϵ -optimal Markov families consisting of a finite number of strategies. The number of strategies depends only on the number of times that the strategy depends only on the current state and the result holds for all finite state gambling problems with a finite state space and a permutation invariant transition probability α from Γ to each element of a finite set Ω .